

Pricing Early-exercise options

GPU Acceleration of SGBM method

Delft University of Technology - Centrum Wiskunde & Informatica

Álvaro Leitao Rodríguez and Cornelis W. Oosterlee
Lausanne - December 4, 2016

Outline

- 1 Definitions
- 2 Basket Bermudan Options
- 3 Stochastic Grid Bundling Method
- 4 Parallel GPU Implementation
- 5 Results
- 6 Conclusions

Definitions

Option

A contract that offers the buyer the right, but not the obligation, to buy (call) or sell (put) a financial asset at an agreed-upon price (the strike price) during a certain period of time or on a specific date (exercise date). Investopedia.

Option price

The fair value to enter in the option contract. In other words, the (discounted) expected value of the contract.

$$V_t = D_t \mathbb{E} [f(S_t)]$$

where f is the *payoff* function, S the underlying asset, t the exercise time and D_t the discount factor.

Definitions - cont.

Pricing techniques

- Stochastic process, S_t .
- Simulation: Monte Carlo method.
- PDEs: Feynman-Kac theorem.

Types of options - Exercise time

- European: End of the contract, $t = T$.
- American: Anytime, $t \in [0, T]$.
- Bermudan: Some predefined times, $t \in \{t_1, \dots, t_M\}$
- Many others: Asian, barrier, ...

Definitions - cont.

Early-exercise option price

- American:

$$V_t = \sup_{t \in [0, T]} D_t \mathbb{E} [f(S_t)].$$

- Bermudan:

$$V_t = \sup_{t \in \{t_1, \dots, t_M\}} D_t \mathbb{E} [f(S_t)].$$

Pricing early-exercise options

- PDEs: Hamilton-Jacobi-Bellman equation.
- Simulation:
 - Least-squares method (LSM), Longstaff and Schwartz.
 - Stochastic Grid Bundling method (SGBM) [JO15].

Basket Bermudan Options

- Right to exercise at a set of times:
 $t \in \{t_0 = 0, \dots, t_m, \dots, t_M = T\}$.
- d -dimensional underlying process: $\mathbf{S}_t = (S_t^1, \dots, S_t^d) \in \mathbb{R}^d$.
- Intrinsic value of the option: $h_t := h(\mathbf{S}_t)$.
- The value of the option at the terminal time T :

$$V_T(\mathbf{S}_T) = f(\mathbf{S}_T) = \max(h(\mathbf{S}_T), 0).$$

- The conditional continuation value Q_{t_m} , i.e. the discounted expected payoff at time t_m :

$$Q_{t_m}(\mathbf{S}_{t_m}) = D_{t_m} \mathbb{E} [V_{t_{m+1}}(\mathbf{S}_{t_{m+1}}) | \mathbf{S}_{t_m}].$$

- The Bermudan option value at time t_m and state \mathbf{S}_{t_m} :

$$V_{t_m}(\mathbf{S}_{t_m}) = f(\mathbf{S}_T) = \max(h(\mathbf{S}_{t_m}), Q_{t_m}(\mathbf{S}_{t_m})).$$

- Value of the option at the initial state \mathbf{S}_{t_0} , i.e. $V_{t_0}(\mathbf{S}_{t_0})$.

Basket Bermudan options scheme

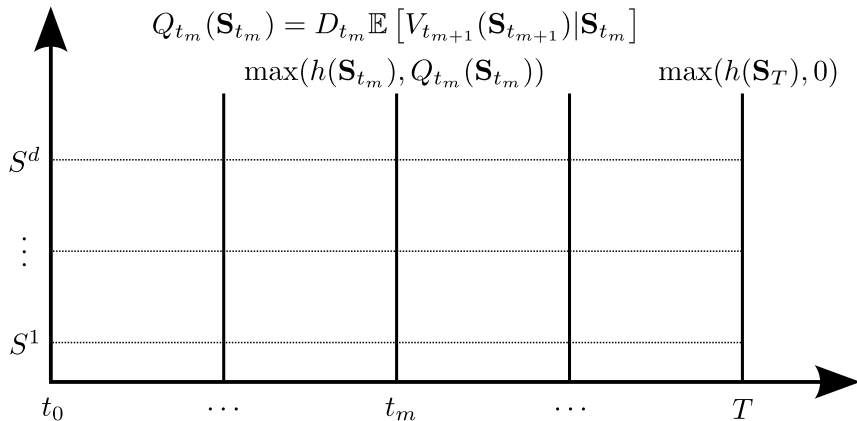


Figure: d-dimensional Bermudan option

Stochastic Grid Bundling Method

- Dynamic programming approach.
- Simulation and regression-based method.
- Forward in time: Monte Carlo simulation.
- Backward in time: Early-exercise policy computation.
- Step I: Generation of stochastic grid points

$$\{\mathbf{S}_{t_0}(n), \dots, \mathbf{S}_{t_M}(n)\}, \quad n = 1, \dots, N.$$

- Step II: Option value at terminal time $t_M = T$

$$V_{t_M}(\mathbf{S}_{t_M}) = \max(h(\mathbf{S}_{t_M}), 0).$$

Stochastic Grid Bundling Method (II)

- Backward in time, t_m , $m \leq M$;
- Step III: Bundling into ν non-overlapping sets or partitions

$$\mathcal{B}_{t_{m-1}}(1), \dots, \mathcal{B}_{t_{m-1}}(\nu)$$

- Step IV: Parameterizing the option values

$$Z(\mathbf{S}_{t_m}, \alpha_{t_m}^\beta) \approx V_{t_m}(\mathbf{S}_{t_m}).$$

- Step V: Computing the continuation and option values at t_{m-1}

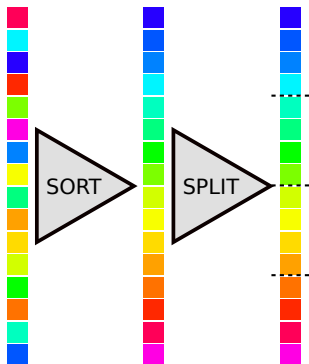
$$\hat{Q}_{t_{m-1}}(\mathbf{S}_{t_{m-1}}(n)) = \mathbb{E}[Z(\mathbf{S}_{t_m}, \alpha_{t_m}^\beta) | \mathbf{S}_{t_{m-1}}(n)].$$

The option value is then given by:

$$\hat{V}_{t_{m-1}}(\mathbf{S}_{t_{m-1}}(n)) = \max(h(\mathbf{S}_{t_{m-1}}(n)), \hat{Q}_{t_{m-1}}(\mathbf{S}_{t_{m-1}}(n))).$$

Bundling

- Original: Iterative process (K-means clustering).
- Problems: Too expensive (time and memory) and distribution.
- New technique: Equal-partitioning. Efficient for parallelization.
- Two stages: sorting and splitting.



Parametrizing the option value

- Basis functions $\phi_1, \phi_2, \dots, \phi_K$.
- In our case, $Z(\mathbf{S}_{t_m}, \alpha_{t_m}^\beta)$ depends on \mathbf{S}_{t_m} only through $\phi_k(\mathbf{S}_{t_m})$:

$$Z(\mathbf{S}_{t_m}, \alpha_{t_m}^\beta) = \sum_{k=1}^K \alpha_{t_m}^\beta(k) \phi_k(\mathbf{S}_{t_m}).$$

- Computation of $\alpha_{t_m}^\beta$ (or $\hat{\alpha}_{t_m}^\beta$) by least squares regression.
- The $\alpha_{t_m}^\beta$ determines the early-exercise policy.
- The continuation value:

$$\begin{aligned} \hat{Q}_{t_{m-1}}(\mathbf{S}_{t_{m-1}}(n)) &= D_{t_{m-1}} \mathbb{E} \left[\left(\sum_{k=1}^K \hat{\alpha}_{t_m}^\beta(k) \phi_k(\mathbf{S}_{t_m}) \right) \mid \mathbf{S}_{t_{m-1}} \right] \\ &= D_{t_{m-1}} \sum_{k=1}^K \hat{\alpha}_{t_m}^\beta(k) \mathbb{E} [\phi_k(\mathbf{S}_{t_m}) \mid \mathbf{S}_{t_{m-1}}]. \end{aligned}$$

Basis functions

- Choosing ϕ_k : the expectations $\mathbb{E} [\phi_k(\mathbf{S}_{t_m}) | \mathbf{S}_{t_{m-1}}]$ should be easy to calculate.
- The intrinsic value of the option, $h(\cdot)$, is usually an important and useful basis function. For example:
 - Geometric basket Bermudan:

$$h(\mathbf{S}_t) = \left(\prod_{\delta=1}^d S_t^\delta \right)^{\frac{1}{d}}$$

- Arithmetic basket Bermudan:

$$h(\mathbf{S}_t) = \frac{1}{d} \sum_{\delta=1}^d S_t^\delta$$

- For \mathbf{S}_t following a GBM: expectations analytically available.

Estimating the option value

- SGBM has been developed as *duality-based method*.
- Provide two estimators (confidence interval).
- *Direct estimator* (high-biased estimation):

$$\widehat{V}_{t_{m-1}}(\mathbf{S}_{t_{m-1}}(n)) = \max \left(h(\mathbf{S}_{t_{m-1}}(n)), \widehat{Q}_{t_{m-1}}(\mathbf{S}_{t_{m-1}}(n)) \right),$$

$$\mathbb{E}[\widehat{V}_{t_0}(\mathbf{S}_{t_0})] = \frac{1}{N} \sum_{n=1}^N \widehat{V}_{t_0}(\mathbf{S}_{t_0}(n)).$$

- *Path estimator* (low-biased estimation):

$$\widehat{\tau}^*(\mathbf{S}(n)) = \min \{ t_m : h(\mathbf{S}_{t_m}(n)) \geq \widehat{Q}_{t_m}(\mathbf{S}_{t_m}(n)), m = 1, \dots, M \},$$
$$v(n) = h(\mathbf{S}_{\widehat{\tau}^*(\mathbf{S}(n))}),$$

$$\underline{V}_{t_0}(\mathbf{S}_{t_0}) = \lim_{N_L} \frac{1}{N_L} \sum_{n=1}^{N_L} v(n).$$

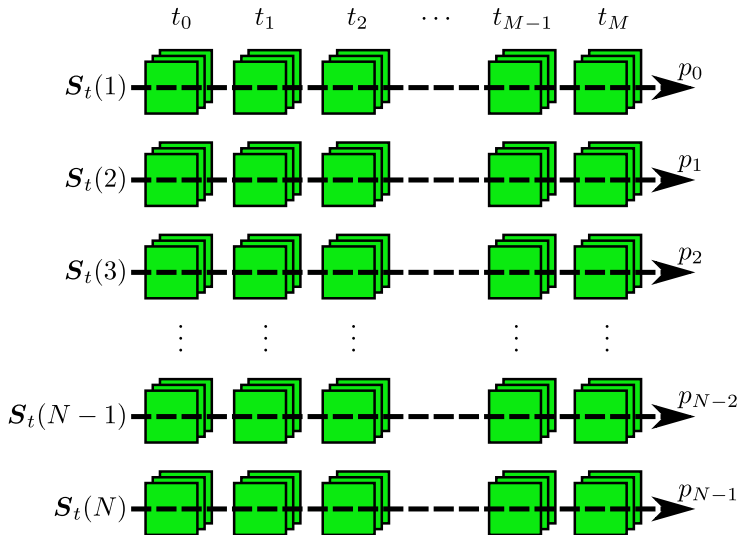
Parallel SGBM on GPU

- NVIDIA CUDA platform.
- Parallel strategy: two parallelization stages:
 - Forward: Monte Carlo simulation.
 - Backward: Bundles → Opportunity of parallelization.
- Novelty in early-exercise option pricing methods.
- Two implementations → K-means vs. Equal-partitioning:
 - K-means: sequential parts.
 - K-means: transfers between CPU and GPU cannot be avoided.
 - K-means: all data need to be stored.
 - K-means: Load-balancing.
 - Equal-partitioning: fully parallelizable.
 - Equal-partitioning: No transfers.
 - Equal-partitioning: efficient memory use.

Parallel SGBM on GPU - Forward in time

- One GPU thread per Monte Carlo simulation.
- Random numbers “on the fly”: cuRAND library.
- Compute intermediate results:
 - Expectations.
 - Intrinsic value of the option.
 - Equal-partitioning: sorting criterion calculations.
- Intermediate results in the registers: fast memory access.
- Original bundling: all the data still necessary.

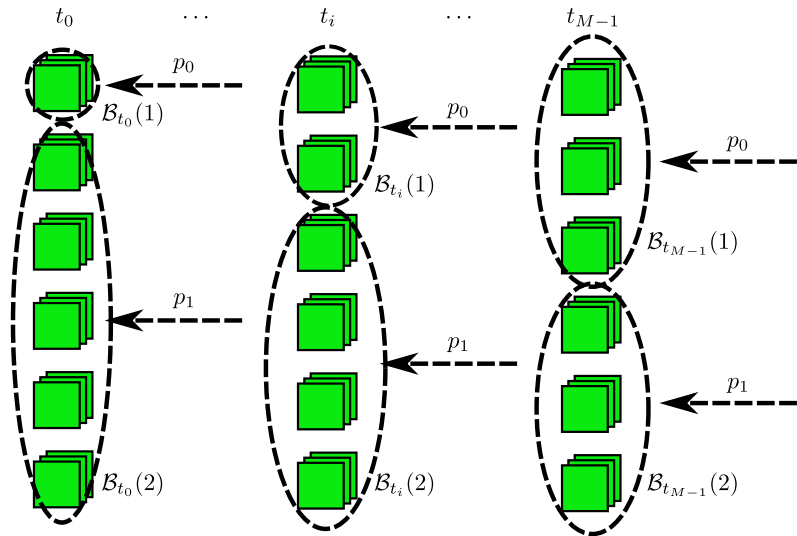
Parallel SGBM on GPU - Forward in time



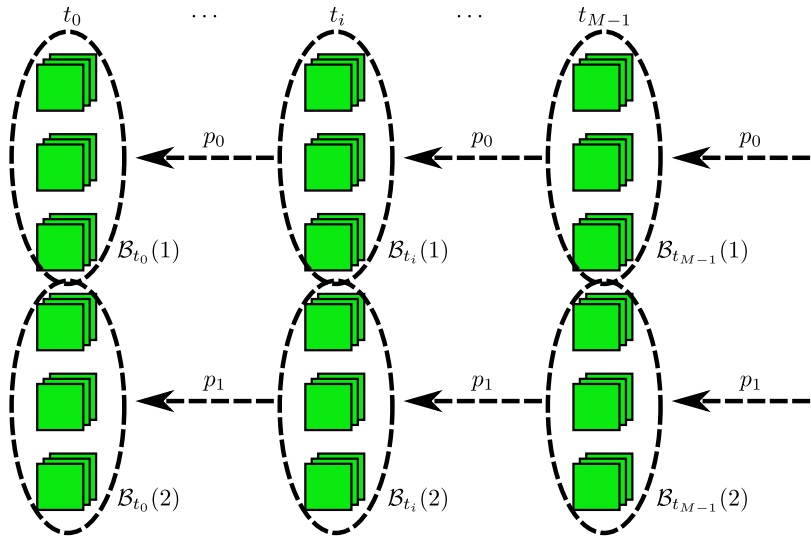
Parallel SGBM on GPU - Backward in time

- One parallelization stage per exercise time step.
- Sort w.r.t bundles: efficient memory access.
- Parallelization in bundles.
- Each bundle calculations (option value and early-exercise policy) in parallel.
- All GPU threads collaborate in order to compute the continuation value.

Parallel SGBM on GPU - Backward in time



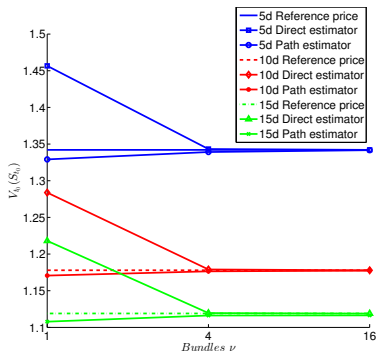
Parallel SGBM on GPU - Backward in time



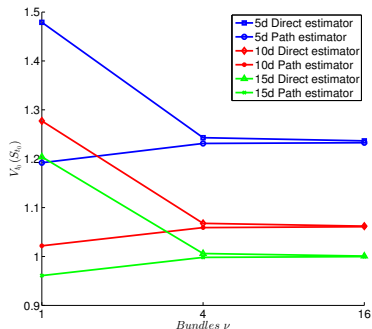
Results

- Accelerator Island system of Cartesius Supercomputer.
 - Intel Xeon E5-2450 v2.
 - NVIDIA Tesla K40m.
 - C-compiler: GCC 4.4.7.
 - CUDA version: 5.5.
- Geometric and arithmetic basket Bermudan put options:
 $\mathbf{S}_{t_0} = (40, \dots, 40) \in \mathbb{R}^d$, $X = 40$, $r_t = 0.06$,
 $\sigma = (0.2, \dots, 0.2) \in \mathbb{R}^d$, $\rho_{ij} = 0.25$, $T = 1$ and $M = 10$.
- Basis functions: $K = 3$.
- Multi-dimensional Geometric Brownian Motion.
- Euler discretization, $\delta t = T/M$.

Equal-partitioning: convergence test



(a) Geometric basket put option



(b) Arithmetic basket put option

Figure: Convergence with equal-partitioning bundling technique. Test configuration: $N = 2^{18}$ and $\Delta t = T/M$.

Speedup

| Geometric basket Bermudan option | | | | | | |
|-----------------------------------------|---------|----------|----------|--------------------|----------|----------|
| | k-means | | | equal-partitioning | | |
| | $d = 5$ | $d = 10$ | $d = 15$ | $d = 5$ | $d = 10$ | $d = 15$ |
| C | 604.13 | 1155.63 | 1718.36 | 303.26 | 501.99 | 716.57 |
| CUDA | 35.26 | 112.70 | 259.03 | 8.29 | 9.28 | 10.14 |
| Speedup | 17.13 | 10.25 | 6.63 | 36.58 | 54.09 | 70.67 |

| Arithmetic basket Bermudan option | | | | | | |
|------------------------------------------|---------|----------|----------|--------------------|----------|----------|
| | k-means | | | equal-partitioning | | |
| | $d = 5$ | $d = 10$ | $d = 15$ | $d = 5$ | $d = 10$ | $d = 15$ |
| C | 591.91 | 1332.68 | 2236.93 | 256.05 | 600.09 | 1143.06 |
| CUDA | 34.62 | 126.69 | 263.62 | 8.02 | 11.23 | 15.73 |
| Speedup | 17.10 | 10.52 | 8.48 | 31.93 | 53.44 | 72.67 |

Table: SGBM total time (s) for the C and CUDA versions. Test configuration: $N = 2^{22}$, $\Delta t = T/M$ and $\nu = 2^{10}$.

Speedup - High dimensions

| Geometric basket Bermudan option | | | | | | |
|-----------------------------------------|----------------|----------|----------|----------------|----------|----------|
| | $\nu = 2^{10}$ | | | $\nu = 2^{14}$ | | |
| | $d = 30$ | $d = 40$ | $d = 50$ | $d = 30$ | $d = 40$ | $d = 50$ |
| C | 337.61 | 476.16 | 620.11 | 337.06 | 475.12 | 618.98 |
| CUDA | 4.65 | 6.18 | 8.08 | 4.71 | 6.26 | 8.16 |
| Speedup | 72.60 | 77.05 | 76.75 | 71.56 | 75.90 | 75.85 |

| Arithmetic basket Bermudan option | | | | | | |
|------------------------------------------|----------------|----------|----------|----------------|----------|----------|
| | $\nu = 2^{10}$ | | | $\nu = 2^{14}$ | | |
| | $d = 30$ | $d = 40$ | $d = 50$ | $d = 30$ | $d = 40$ | $d = 50$ |
| C | 993.96 | 1723.79 | 2631.95 | 992.29 | 1724.60 | 2631.43 |
| CUDA | 11.14 | 17.88 | 26.99 | 11.20 | 17.94 | 27.07 |
| Speedup | 89.22 | 96.41 | 97.51 | 88.60 | 96.13 | 97.21 |

Table: SGBM total time (s) for a high-dimensional problem with equal-partitioning. Test configuration: $N = 2^{20}$ and $\Delta t = T/M$.

Conclusions

- Efficient parallel GPU implementation.
- Extend the SGBM's applicability: Increasing dimensionality.
- New bundling technique.
- Future work:
 - Explore the new CUDA 7 features: cuSOLVER (QR factorization).
 - CVA calculations.

References



Shashi Jain and Cornelis W. Oosterlee.

The Stochastic Grid Bundling Method: Efficient pricing of Bermudan options and their Greeks.

Applied Mathematics and Computation, 269:412–431, 2015.



Álvaro Leitao and Cornelis W. Oosterlee.

GPU Acceleration of the Stochastic Grid Bundling Method for Early-Exercise options.

International Journal of Computer Mathematics,
92(12):2433–2454, 2015.

Acknowledgements



Thank you for your attention

Appendix

- Geo. basket Bermudan option - Basis functions:

$$\phi_k(\mathbf{S}_{t_m}) = \left(\left(\prod_{\delta=1}^d S_{t_m}^{\delta} \right)^{\frac{1}{d}} \right)^{k-1}, \quad k = 1, \dots, K,$$

- The expectation can directly be computed as:

$$\mathbb{E} [\phi_k(\mathbf{S}_{t_m}) | \mathbf{S}_{t_{m-1}}(n)] = \left(P_{t_{m-1}}(n) e^{\left(\bar{\mu} + \frac{(k-1)\bar{\sigma}^2}{2} \right) \Delta t} \right)^{k-1},$$

where,

$$P_{t_{m-1}}(n) = \left(\prod_{\delta=1}^d S_{t_{m-1}}^{\delta}(n) \right)^{\frac{1}{d}}, \quad \bar{\mu} = \frac{1}{d} \sum_{\delta=1}^d \left(r - q_{\delta} - \frac{\sigma_{\delta}^2}{2} \right), \quad \bar{\sigma}^2 = \frac{1}{d^2} \sum_{p=1}^d \left(\sum_{q=1}^d C_{pq}^2 \right)^2$$

Appendix

- Arith. basket Bermudan option - Basis functions:

$$\phi_k(\mathbf{S}_{t_m}) = \left(\frac{1}{d} \sum_{\delta=1}^d S_{t_m}^{\delta} \right)^{k-1}, \quad k = 1, \dots, K.,$$

- The summation can be expressed as a linear combination of the products:

$$\left(\sum_{\delta=1}^d S_{t_m}^{\delta} \right)^k = \sum_{k_1+k_2+\dots+k_d=k} \binom{k}{k_1, k_2, \dots, k_d} \prod_{1 \leq \delta \leq d} \left(S_{t_m}^{\delta} \right)^{k_{\delta}},$$

- And the expression for Geometric basket option can be applied.