## Quantum computing for computational finance

overview, challenges and opportunities

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## Motivation

- Quantum computers could bring unparalleled competitive advantage to financial companies in areas like portfolio optimisation, option pricing, quantitative risk management or Machine Learning models.
- Quantum computers are able to handle exponentially growing (in qubits) Hilbert spaces.
- Thus, quantum computing becomes an attractive framework for calculations over large multi-dimensional domains.
- Quantum algorithms could potentially overcome their classical counterparts in dealing with combinatorial explosions and the curse of dimensionality.
- However, bringing this to practice encounters several bottlenecks, especially with the current or near-term quantum technologies (NISQ).
- Then, plenty of room for contributions!


## Disclaimer

- Quantum computing literature is experiencing an explosion: This overview incorporates only a few of the current trends.
- The selection of the addressed topics reflects only my view (interests) within the vast scope of the computational finance field.
- Then, many important topics are not addressed here: optimal investment, time series, blockchain, cryptography, etc.
- There might be inconsistencies or certain abuse in the (mathematical and/or quantum) notation. In some cases, that is intentional, for the sake of clarity. In others...sorry in advance!


## Outline

## Quantum Computing basics

Overview
Quantum Monte Carlo
Quantum Financial PDEs
Quantum Machine/Deep Learning

Challenges

Opportunities

# Quantum Computing basics 

## Quantum Computing basics (I)

- The basic unit of information is the qubit (alternatively to the bit).
- A qubit is represented by a (column) vector:

$$
\binom{\alpha}{\beta}
$$

with the amplitudes $\alpha, \beta \in \mathbb{C}$ and $|\alpha|^{2}+|\beta|^{2}=1$.

- Basis states:

$$
|0\rangle=\binom{1}{0}, \quad|1\rangle=\binom{0}{1}
$$

- $\{|0\rangle,|1\rangle\}$ is a computational basis for a quantum state:

$$
|\psi\rangle=\alpha|0\rangle+\beta|1\rangle
$$

- When measuring the state:
- get 0 with probability $|\alpha|^{2}$
- get 1 with probability $|\beta|^{2}$


## Quantum Computing basics (II)

- The Bloch sphere provides a representation of qubit state
- Measuring a qubit occurs along the $Z$ axis, so it is irreversible and will collapse to either 0 or 1



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## Quantum Computing basics (III)

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- Gates are reversible and can be represented as unitary matrices acting on the qubit vectors.



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## Quantum Computing basics (IV)

- Superposition: Identically prepared qubits can still behave randomly
- The randomness is inherent in the quantum nature

|1

$\sim 50 / 50$ chance of being $|0\rangle$ or $|1\rangle$


## Quantum Computing basics (and V)

- Each row represents a bit, either quantum or classical
- The operations are performed each qubit from left to right
- Measurement to extract the information



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Overview

## Overview

- Based on the publication [11]:
A. Gómez, Á. Leitao, A. Manzano, D. Musso, M.R. Nogueiras, G.

Ordóñez and C. Vázquez. A survey on quantum computational finance for derivatives pricing and VAR, Archives of
Computational Methods in Engineering 29: 4137-4163, 2022.

- Survey on the classical methods for pricing and VaR, and their potential quantum counterparts.
- We mainly focus on:
- Monte Carlo-like methods.
- Partial differential equations (PDEs).
- Machine Learning/Neural Networks/Deep Learning.


## Quantum Monte Carlo

## Quantum Monte Carlo (QMC)

The quantum-accelarated Monte Carlo could potentially/theoretically provide a quadratic speedup.

How? Quantum Amplitude Estimation.
Monte Carlo methods (for pricing) can be informally defined as

$$
\frac{1}{M} \sum_{i=0}^{M-1} f\left(X_{i}\right) \approx \mathbb{E}[f(X)]=\int f(x) p(x) d x \approx \sum_{j=0}^{N-1} f\left(x_{j}\right) p\left(x_{j}\right)
$$

where $p(x)$ is a density function.
Analogously, Quantum Monte Carlo (QMC) assumes a state of the form

$$
|\psi\rangle=|0\rangle \otimes \sum_{j=0}^{N-1} f\left(x_{j}\right) p\left(x_{j}\right)|j\rangle+|1\rangle \otimes \sum_{j=0}^{N-1} \sqrt{1-f^{2}\left(x_{j}\right) p\left(x_{j}\right)}|j\rangle
$$

where the quantity of interest is encoded in the state's amplitude.

## QMC: Quantum Amplitude Estimation

Given a state:

$$
|\psi\rangle=a|\phi\rangle+\sqrt{1-a^{2}}\left|\phi^{\perp}\right\rangle
$$

Quantum Amplitude Estimation (QAE) is an algorithm which gives an estimation $\hat{a} \pm \frac{\epsilon}{2}$ of the amplitude $a$.


This technique promises to obtain a quadratic speedup over its classical counterpart.

To achieve so it relies on two main subroutines:

- Grover (search) amplification.
- Quantum Phase Estimation.


## QMC: Grover Amplification

Given a state:

$$
|\psi\rangle=\sin (\theta)|\phi\rangle+\cos (\theta)\left|\phi^{\perp}\right\rangle
$$

Grover operator performs the following transformation:

$$
\mathcal{Q}^{k}|\psi\rangle=\sin ((2 k+1) \theta)|\phi\rangle+\cos ((2 k+1) \theta)\left|\phi^{\perp}\right\rangle .
$$



## QMC: Quantum Amplitude Estimation (graphically)



## QMC: Quantum Amplitude Estimation (circuit)



Figure 1: Quantum Amplitude Estimation

## QMC: Quantum Amplitude Estimation (convergence)

Theorem (Mean estimation for [ 0,1 ] bounded functions [25]) Let there be given a quantum circuit $\mathcal{P}$ on $n$ qubits. Let $v(\mathcal{P})$ be the random variable that maps to $v(x) \in[0,1]$ when the bit string $x$ is measured as the output of $\mathcal{P}$. Let $\mathcal{R}$ be defined as

$$
\mathcal{R}|x\rangle|0\rangle=|x\rangle(\sqrt{1-v(x)}|0\rangle-\sqrt{v(x)}|1\rangle) .
$$

Let $|\mathcal{X}\rangle$ be defined as $|\mathcal{X}\rangle=\mathcal{R}\left(\mathcal{P} \otimes \mathcal{I}_{2}\right)\left|0^{n+1}\right\rangle$. Set $\mathcal{U}=\mathcal{I}_{2^{n+1}}-2|\mathcal{X}\rangle\langle\mathcal{X}|$. There exists a quantum algorithm that uses $\mathcal{O}(\log 1 / \delta)$ copies of the state $\mathcal{X}$, uses $\mathcal{U}$ for a number of times proportional to $\mathcal{O}(m \log 1 / \delta)$ and outputs an estimate $\hat{\mu}$ such that

$$
|\hat{\mu}-\mathbb{E}[v(\mathcal{P}])| \leq C\left(\frac{\sqrt{\mathbb{E}[v(\mathcal{P})]}}{m}+\frac{1}{m^{2}}\right)
$$

with probability at least $1-\delta$, where $C$ is a universal constant. In particular, for any fixed $\delta>0$ and any $\epsilon$ such that $0<\epsilon \leq 1$, to produce an estimate $\hat{\mu}$ such that, with probability at least $1-\delta,|\hat{\mu}-\mathbb{E}[v(\mathcal{P})]| \leq \epsilon \mathbb{E}[v(\mathcal{P})]$, it suffices to take $m=\mathcal{O}\left((\epsilon \mathbb{E}[v(\mathcal{P})])^{-1}\right)$. To achieve $|\hat{\mu}-\mathbb{E}[v(\mathcal{P})]| \leq \epsilon$ with probability at least $1-\delta$, it suffices to take $m=\mathcal{O}\left(\epsilon^{-1}\right)$.

## QMC: variations on QAE

- Quantum Amplitude Estimation is not feasible with the current technology as it depends on Quantum Phase Estimation.
- The depth of the circuit due to the use of a Quantum Fourier Transform (QFT) is prohibitive.
- To mitigate it several new techniques have appeared:
- Simplified Quantum Counting (SQAE)[1]
- Maximum Likelihood Amplitude Estimation (MLAE)[34, 35]
- Iterative Quantum Amplitude Estimation (IQAE)[14]
- We have proposed another alternative [23]:
A. Manzano, D. Musso, Á. Leitao. Real Quantum Amplitude Estimation, EPJ Quantum Technology 10(2), 2023.


## RQAE: Settings

- Consider a one-parameter family of oracles $\mathcal{A}_{b}$ that, acting on the state $|0\rangle$, yield

$$
\mathcal{A}_{b}|0\rangle=|\psi\rangle=(a+b)|\phi\rangle+c_{b}\left|\phi^{\perp}\right\rangle_{b},
$$

where $a$ is a real number, $b$ is an auxiliary, continuous and real parameter that we call "shift".

- Given a precision level $\epsilon$ and a confidence level $1-\gamma$, the goal of the RQAE algorithm is to compute an interval $\left(a_{l}^{\min }, a_{l}^{\max }\right) \subset[-1,1]$ of width smaller than $2 \epsilon$ which contains the value of a with probability greater or equal to $1-\gamma$.



## RQAE: Algorithm (graphically)



## RQAE: Main properties

Given $\epsilon, \gamma$ and an amplification policy $\left(k_{i}, i=0, \ldots, l\right)$ :

- Circuit depth bounded by:

$$
k_{I} \leq\left\lceil\frac{\arcsin \left(\sqrt{2 \epsilon^{\rho}}\right)}{\arcsin (2 \epsilon)}-\frac{1}{2}\right\rceil=k^{\max } .
$$

- Precision $\epsilon$ with confidence $1-\gamma$ (Proof of Correctness):

$$
\mathbb{P}\left[a \notin\left(a_{l}^{\min }, a_{l}^{\max }\right)\right] \leq \gamma .
$$

- The total number of calls to the oracle is bounded by:

$$
N_{\text {oracle }}<C_{1} \frac{1}{\epsilon} \log \left(\frac{C_{2}}{\gamma}\right)
$$

where the constants $C_{1}$ and $C_{2}$ depend on the amplification policy.

## RQAE: Numerical results (I)



Figure 2: Number of calls to the oracle $N_{\text {oracle }}$ versus the required precision $\epsilon$.

## RQAE: Numerical results (and II)



## RQAE: Main Conclusions

- More information from the quantum circuit than just the module of the amplitude, i.e., the sign of the quantity of interest, increasing the applicability.
- RQAE is an iterative algorithm which offers explicit control over the amplification policy through adjustable parameters.
- Control (also via the free parameters) the depth of the circuit, a crucial feature in the current NISQ era.
- A rigorous (and clean) theoretical analysis of the RQAE performance is provided, proving that it achieves a quadratic speedup (w.r.t. unamplified sampling), modulo logarithmic corrections.


## Quantum Financial PDEs

## Quantum approaches for Black-Scholes PDE

- Some financial PDEs can be mapped into the propagation governed by an appropriate Hamiltonian operator [12, 9].
- Applying the change of variable $S=e^{x}$ on the Black-Scholes eq.,

$$
\frac{\partial V}{\partial t}+\left(\mu-\frac{\sigma^{2}}{2}\right) \frac{\partial V}{\partial x}+\frac{\sigma^{2}}{2} \frac{\partial^{2} V}{\partial x^{2}}-\mu V=0
$$

which can be written as a Schrödinger-like equation,

$$
\frac{\partial V}{\partial t}=-i \hat{H}_{\mathrm{BS}} V
$$

where

$$
\hat{H}_{\mathrm{BS}}=i \frac{\sigma^{2}}{2} \hat{p}^{2}-\left(\frac{\sigma^{2}}{2}-\mu\right) \hat{p}+i \mu \mathbb{I}, \quad \text { with } \quad \hat{p}=-i \frac{\partial}{\partial x} .
$$

- The Hamiltonian $\hat{H}_{\mathrm{BS}}$ is not Hermitian.
- Therefore, the associated evolution operator $\hat{U}\left(t, t_{0}\right)=e^{-i \hat{H}_{\mathrm{BS}}\left(t-t_{0}\right)}$ is not unitary.


## PDEs: Real-time propagation

- To implement $\hat{U}\left(t, t_{0}\right)$ into a quantum circuit, one can consider an enlarged system, i.e. a doubled unitary operator [12].
- Require of adding an auxiliary qubit.
- $\hat{H}_{\mathrm{BS}}$ is diagonal in momentum space $\rightarrow$ diagonal operator $\rightarrow$ QFT (and Inverse QFT) $\rightarrow$ exponential speedup.
- But, an overall exponential speedup requires efficient loading of the model and payoff function.
- Again, QFT (IQFT) is gate-wise demanding (incompatible NISQ).


Measure and postselection

## PDEs: Real-time propagation (solution)

- The algorithm achieves a high degree of agreement in a fault-tolerant quantum computer...
- ...but with a $60 \%$ success probability in the measurement and post-selection (depending on the financial parameters).
- Not tested in a real NISQ quantum system.

(a) Boundary error without duplication.

(b) Convergence in qubits (point).


## PDEs: Imaginary-time propagation

- Additional change of variable $\tau=\sigma^{2}(T-t)$ and transformation $v(x, \tau)=\exp (-a x-b \tau) V(t, s)$, with suitable constants $a$ and $b$,

$$
\frac{\partial v}{\partial \tau}=\frac{1}{2} \frac{\partial^{2} v}{\partial x^{2}} .
$$

- Using the Wick rotation $\tilde{\tau}=-i \tau$ (real time to imaginary time), the heat equation turns into a Schrödinger-like equation,

$$
\frac{\partial v}{\partial \tilde{\tau}}=-\hat{H}_{H E} v,
$$

where

$$
\hat{H}_{\mathrm{HE}}=-\frac{i}{2} \hat{q}^{2}, \quad \text { with } \quad \hat{q}=-i \frac{\partial}{\partial x} .
$$

- This leads to a purely anti-Hermitian Hamiltonian operator.
- Imaginary-time propagation transforms oscillations into dampings.
- Problem of finding the ground state of quantum systems, well investigated in condensed matter physics and chemistry.
- The imaginary time evolution operator is approximated by an ansatz circuit in [9].


## PDEs: Imaginary-time propagation (solution)

- The solution is retrieved by a hybrid quantum-classical algorithm:



(c) Prices of European option.
(d) Errors of European option.

Quantum Machine/Deep Learning

## Quantum Machine/Deep Learning

- (Quantum) Principal Component Analysis:
- Eigenvalues by Quantum Phase Estimation [27].
- Convert covariance matrix into density matrix (QPCA) [20, 2].
- (Quantum) Regression:
- Solving linear systems by the HHL algorithm [37].
- Quantum Kernel Estimation [8].
- Quantum regression with Gaussian processes [38].
- Hybrid classical-quantum deep learning:
- Move the training to a quantum computer (quantum annealing) [3].
- Quantum-enhanced reinforcement learning [29].
- Quantum GANs [26].
- Boltzman machines $\rightarrow$ Born machines [36, 4].
- Full Quantum Neural Network (QNN):
- NN models based on the principles of quantum mechanics [17].
- How to train QNN? Recent advances in [6, 7].
- Promising approach: Parametrized Quantum Circuits


## Parametrized Quantum Circuits (PQCs)

- Also known as variational circuits or quantum circuit learning.
- First theoretical results on accessibility, expressivity and universality.
- Circuits with both fixed and adjustable ("parametrized") gates.
- The training is carried out by a classical optimiser.
- Each layer composed by a trainable circuit block $W_{i}(\theta)$ and a data-encoding block $S(x)$ :


Figure 3: Parametrized Quantum Circuit.

## PQCs: trigonometric series

- A PQC model can be written as a generalized trigonometric series:

$$
\mathbb{E}[M]=\langle 0| U^{\dagger}(\boldsymbol{x} ; \boldsymbol{\theta}) M U(\boldsymbol{x} ; \boldsymbol{\theta})|0\rangle=f(\boldsymbol{x} ; \boldsymbol{\theta})=\sum_{\boldsymbol{\omega} \in \Omega} c_{\omega}(\theta) e^{i \boldsymbol{\omega} x},
$$

where $M$ is an observable, $U(\boldsymbol{x} ; \boldsymbol{\theta})$ is a quantum circuit that depends on inputs $\boldsymbol{x}=\left(x_{0}, x_{1}, \ldots, x_{N}\right)$ and the parameters $\boldsymbol{\theta}=\left(\theta_{0}, \theta_{1}, \ldots, \theta_{T}\right)$.

- Accessibility: with $\Omega \subset \mathbb{Z}^{N} \rightarrow$ (partial) Fourier series!
- The coefficients $c_{\omega}$ determine the expressivity (how the accessible functions can be combined).
- But the expressivity is also limited by the data encoding strategy.
- Universality: the Fourier series formalism allows to study quantum models using the results in Fourier analysis (see [30] and [21]).


## PQCs: Universality results (I)

## Definition

Let $U(\boldsymbol{x} ; \boldsymbol{\theta})$ be modelled as a unitary such that (1 layer):

$$
U(\boldsymbol{\theta}, \boldsymbol{x})=W^{(2)}\left(\boldsymbol{\theta}^{(2)}\right) S(\boldsymbol{x}) W^{(1)}\left(\boldsymbol{\theta}^{(1)}\right),
$$

and

$$
S(\boldsymbol{x})=e^{-x_{1} H} \otimes \cdots \otimes e^{-x_{N} H}=: S_{H}(\boldsymbol{x})
$$

where $H$ is a particular Hamiltonian.

## Definition

Let $\left\{H_{m} \mid m \in \mathbb{N}\right\}$ be a Hamiltonian family where $H_{m}$ acts on $m$ subsystems of dimension $d$. Such a Hamiltonian family gives rise to a family of models $\left\{f_{m}\right\}$ in the following way:

$$
\begin{equation*}
f_{m}(\boldsymbol{x})=\langle\Gamma| S_{H_{m}}^{\dagger}(\boldsymbol{x}) M S_{H_{m}}(\boldsymbol{x})|\Gamma\rangle . \tag{1}
\end{equation*}
$$

with $|\Gamma\rangle:=W^{(1)}\left(\boldsymbol{\theta}^{(1)}\right)|0\rangle$.

## PQCs: Universality results (II)

Theorem (Convergence in $L^{2}$ [30])
Let $\left\{H_{m}\right\}$ be a universal Hamiltonian family, and $\left\{f_{m}\right\}$ the associated quantum model family, defined via (1). For all functions $f^{*} \in L^{2}\left([0,2 \pi]^{N}\right)$, and for all $\epsilon>0$, there exists some $m^{\prime} \in \mathbb{N}$, some state $|\Gamma\rangle \in \mathbb{C}^{m^{\prime}}$ and some observable $M$ such that

$$
\left\|f_{m^{\prime}}-f^{*}\right\|_{L^{2}}<\epsilon .
$$

Theorem (Convergence in $L^{p}$ [21])
Let $\left\{H_{m}\right\}$ be a universal Hamiltonian family, and $\left\{f_{m}\right\}$ the associated quantum model family, defined via (1). For all functions $f^{*} \in L^{p}\left([0,2 \pi]^{N}\right)$ where $1 \leq p<\infty$, and for all $\epsilon>0$, there exists some $m^{\prime} \in \mathbb{N}$, some state $|\Gamma\rangle \in \mathbb{C}^{m^{\prime}}$, and some observable $M$ such that:

$$
\left\|f_{m^{\prime}}-f^{*}\right\|_{L^{p}}<\epsilon .
$$

## PQCs: Universality results (and III)

Theorem (Convergence in $C^{0}$ [21])
Let $\left\{H_{m}\right\}$ be a universal Hamiltonian family, and $\left\{f_{m}\right\}$ the associated quantum model family, defined via (1). For all functions $f^{*} \in C^{0}(U)$ where $U$ is compactly contained in the closed cube $[0,2 \pi]^{N}$, and for all $\epsilon>0$, there exists some $m^{\prime} \in \mathbb{N}$, some state $|\Gamma\rangle \in \mathbb{C}^{m^{\prime}}$, and some observable $M$ such that $f_{m^{\prime}}$ converges uniformly to $f^{*}$ :

$$
\left\|f_{m^{\prime}}-f^{*}\right\|_{C^{0}}<\epsilon
$$

with

$$
\left\|f_{m^{\prime}}-f^{*}\right\|_{C^{0}}:=\sup _{\mathbf{x} \in[0,2 \pi]^{N}}\left\|f_{m^{\prime}}(\mathbf{x})-f^{*}(\mathbf{x})\right\|
$$

## Challenges

## Discussion on Quantum Monte Carlo

Is the Quantum Monte Carlo what we (computational finance community) expect?

In [33] they divide the routine for computing the price of a plain vanilla in three steps:


They promise a quadratic speedup over classical Monte Carlo:
"This represents a theoretical quadratic speed-up compared to classical Monte Carlo methods."

## Classical Monte Carlo vs Quantum Monte Carlo

When claiming a "quadratic" speedup of the QMC over the Classical, what are they comparing?

Steps involved in Classical and Quantum Monte Carlo :

| Quantum Monte Carlo | Classical Monte Carlo |
| :--- | :---: |
| Load Distribution | Load parameters... <br> Simulate the paths |
| Load Payoff | Compute payoff |
| Amplitude Estimation | Sum over paths <br> Print the results |

"In most of the existing literature on option pricing for equities using quantum computers... an SDE is tacitly solved... Once this SDE is solved... the pricing of a particular security begins by applying QAE.[5]"

## Bottleneck

The bottleneck in Classical Monte Carlo is in simulating paths. Analogously, the bottleneck in the quantum algorithm is in the loading/simulation/computation of the distribution.


The quantum advantage might disappear when taking into account the cost of simulation:
"Although preparing such states is in principle always possible for reasonable stochastic processes, efficient realization of this method demands a careful analysis and may not always result in a practical quantum advantage." (see [5])

## Quantum Algorithm Pipeline

Quantum advantage of an algorithm? Efficient end-to-end framework!
We propose (in [22], inspired in [8]) a pipeline as:


- Focus on specific problems $\longrightarrow$ Modularity
- Use algorithms from other works $\longrightarrow$ Reusability


## Other challenges for problems in quantitative finance

- Data loading for Quantum Machine Learning models.
- Quantum-native function implementations (using unitary transforms).
- Information extraction from a quantum state:
- QAE can be seen as an efficient information extraction routine
- Post selection in PDE-Hamiltonian simulation algoritms?
- Rigorous proofs for:
- Speedups (quantum advantage)
- Estimation convergence
- Circuits complexity (depth)
- Quantum volume (NISQ):
- Intrinsic noisy of the current quantum systems (the shallower the better)
- Limited number of qubits (i.e. to represent floating-point numbers)
- Others: coherence time, measurement errors, circuit compiler efficiency, etc.


## Opportunities

## Opportunities

- Quantum Monte Carlo: All the described above!
- PDEs:
- Mapping to linear systems (classical numerical methods)
- Proofs of algorithm complexity (close to exponential speedup?)
- Proofs for the universality/expressivity of quantum DL models:
- Theory behind the general reproducing kernels
- Adapt/use QML algorithms for financial applications.
- Efficient quantum versions of successful classical algorithms:
- Quantum COS
- PINNs
- Differential Machine Learning
- Other:
- Alternatives to Harrow-Hassidim-Lloyd (HHL) for linear systems.
- Similar opportunities in optimization algorithms.
- At practical level: NISQ era!


## Conclusions

- In recent years we have seen significant advances in quantum algorithms with application to financial mathematical problems.
- While this progress is very encouraging, further work will be required to prove that Quantum Computing can deliver real-world advantage.
- Especially if this advantage is to be delivered on NISQ technology with limitations to both the number of logical qubits and the depth of quantum circuits.
- Research into financial applications of quantum computing is accelerating with new ideas emerging at rapid pace...
- ...but important breakthroughs across the technology stack will be needed to make the approaches viable.
- Theory/software is ahead of practice/hardware!


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## QMC: Risk measures

Find $\left.\operatorname{VaR}_{\alpha}(X)=\inf \{x: \mathbb{P}[X \leq x]\} \geq 1-\alpha\right\}=\inf \left\{x: F_{X}(x) \geq 1-\alpha\right\}:$

$$
f_{J}(x)= \begin{cases}1 & \text { if } x \leq x_{J} \\ 0 & \text { otherwise }\end{cases}
$$

Thus, the original QMC state becomes

$$
|\psi\rangle=|0\rangle \otimes \sum_{j=J+1}^{N-1} p\left(x_{j}\right)|j\rangle+|1\rangle \otimes \sum_{j=0}^{J} \sqrt{p\left(x_{j}\right)}|j\rangle
$$

A bisection search over $J$ and measuring $|1\rangle$ gives the $x_{J_{\alpha}} \approx \operatorname{VaR}_{\alpha}(X)$
To estimate $\operatorname{CVaR}_{\alpha}(X)$, take $f(x)=\frac{x}{\chi_{J_{\alpha}}} f_{J_{\alpha}}(x)$, so
$\left.|\psi\rangle=|0\rangle \otimes\left(\sum_{j=J_{\alpha}+1}^{N-1} p\left(x_{j}\right)|j\rangle+\sum_{j=0}^{J_{\alpha}}\left(1-\frac{x_{j}}{x_{J_{\alpha}}}\right) p\left(x_{j}\right)|j\rangle\right)+|1\rangle \otimes \sum_{j=0}^{J_{\alpha}} \sqrt{\frac{x_{j}}{x_{J_{\alpha}}} p\left(x_{j}\right)} \right\rvert\, j$
and measure $|1\rangle$. Then, $\operatorname{CVaR}_{\alpha}(X) \approx \frac{x_{J_{\alpha}}}{1-\alpha} \sum_{j=0}^{J_{\alpha}} \frac{x_{j}}{x_{J_{\alpha}}} p\left(x_{j}\right)$

## QMC: Pricing

Using Y-rotations and a comparator (in $K$ ), we can construct:

$$
\begin{aligned}
|\psi\rangle & =|0\rangle \otimes \sum_{x_{j}<K} \sqrt{p\left(x_{j}\right)}|j\rangle\left[\cos \left(g_{0}\right)|0\rangle+\sin \left(g_{0}\right)|1\rangle\right] \\
& +|1\rangle \otimes \sum_{x_{j} \geq K} \sqrt{p\left(x_{j}\right)}|j\rangle\left[\cos \left(g_{0}+g\left(x_{j}\right)\right)|0\rangle+\sin \left(g_{0}+g\left(x_{j}\right)\right)|1\rangle\right]
\end{aligned}
$$

The probability of measuring the second ancilla (auxiliary) state $|1\rangle$ is:

$$
P=\sum_{x_{j}<K} p\left(x_{j}\right) \sin ^{2}\left(g_{0}\right)+\sum_{x_{j} \geq K} p\left(x_{j}\right) \sin ^{2}\left(g_{0}+g\left(x_{j}\right)\right)
$$

For a European call $\left(\max \left(0, x_{j}-K\right)\right)$, set $g(x)=\frac{2 c(x-K)}{x_{\max }-K}, g_{0}=\frac{\pi}{4}-c$.
Thus, using that $\sin ^{2}\left(c f(x)+\frac{\pi}{4}\right)=c f(x)+\frac{1}{2}+\mathcal{O}\left(c^{3} f^{3}(x)\right)$, we have

$$
\begin{aligned}
P & \approx \sum_{x_{j}<K} p\left(x_{j}\right)\left(\frac{1}{2}-c\right)+\sum_{x_{j} \geq K} p\left(x_{j}\right)\left(\frac{2 c\left(x_{j}-K\right)}{x_{\max }-K}+\frac{1}{2}-c\right) \\
& =\frac{1}{2}-c+\frac{2 c}{x_{\max }-K} \sum_{x_{j} \geq K} p\left(x_{j}\right)\left(x_{j}-K\right)
\end{aligned}
$$

## Algorithms for simulating/loading distributions

Some approaches to compute the initial distribution relies on:

1. Use a general method to convert unitary operators to circuits [13].
2. Use specific methods to initialize the amplitudes to a normalized vector [31, 18, 32].
3. Use the properties of the probability distribution to create an efficient circuit [15, 24].
4. Create an ad-hoc circuit using Parameterized Quantum Circuit (PQC) which approximates the amplitudes [26].
5. Using Tensor Networks techniques [28, 16, 10].

Approaches for financial problems which aim to simulate the underlying SDE:

- Analogous SDE simulation with a "quantum" Floating-Point number representation [25]
- Simulate the MC paths using a trinomial tree[19].
- Use Feynmann-Kac approach[5].


## Quantum Matrix

Instead of working with a quantum state of the form:

$$
|\psi\rangle=\sum_{i=0}^{I-1} \sum_{j=0}^{J-1} c_{i j}|i\rangle_{n_{l}} \otimes|j\rangle_{n \jmath}
$$

we work with a data structure (quantum matrix) storing the same information:

|  | $\|0\rangle_{n_{I}}$ | $\|1\rangle_{n_{I}}$ | $\ldots$ | $\|j\rangle_{n_{I}}$ | $\ldots$ | $\|J\rangle_{n_{I}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\|0\rangle_{n_{I}}$ | $c_{00}$ | $c_{01}$ | $\ldots$ | $c_{0 j}$ | $\ldots$ | $c_{0 J}$ |
| $\|1\rangle_{n_{I}}$ | $c_{10}$ | $c_{11}$ | $\ldots$ | $c_{1 j}$ | $\ldots$ | $c_{1 J}$ |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |
| $\|i\rangle_{n_{I}}$ | $c_{i 0}$ | $c_{i 1}$ | $\ldots$ | $c_{i j}$ | $\ldots$ | $c_{i J}$ |
| $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ | $\ldots$ |
| $\|I\rangle_{n_{I}}$ | $c_{I 0}$ | $c_{I 1}$ | $\ldots$ | $c_{I j}$ | $\ldots$ | $c_{I J}$ |

## Arithmetic on quantum matrices

- We provide arithmetic operations specifically designed for our framework.
- We distinguish three categories depending on their efficiency.
- The most efficient ones (green) depend on the co-dimension instead of the array's length:
- Sum and difference.
- Multiplication by scalar.
- The techniques with middle efficiency (yellow) involve big number of multicontrolled operations but do not depend on an oracle:
- Permutations in the quantum matrix.
- The least efficient techniques (red) depend on an oracle:
- Element-wise squaring of an array.
- Scalar product.


## Advantages \& Disadvantages

The main advantages of the proposed framework are:

- Easy to understand (the pipeline and the quantum matrix/tensor).
- Good level of abstraction (from the applications perspective and the implementation one).
- Efficient performance for operations on large structures.

Nevertheless, some drawbacks may arise for certain tasks:

- Perform operations on single elements $\longrightarrow$ The efficiency depends on the codimension.
- Multiplication of arrays $\longrightarrow$ Unitary operators do not produce multiplication.

